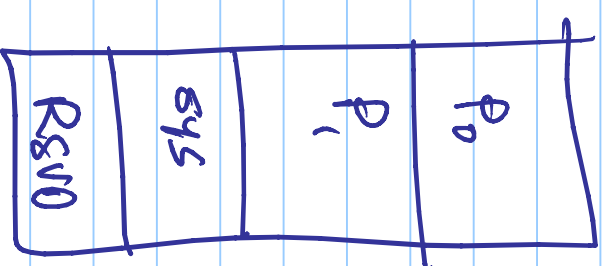
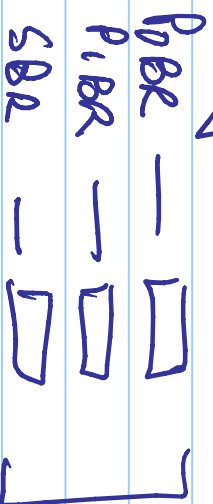


VAX virtual mem

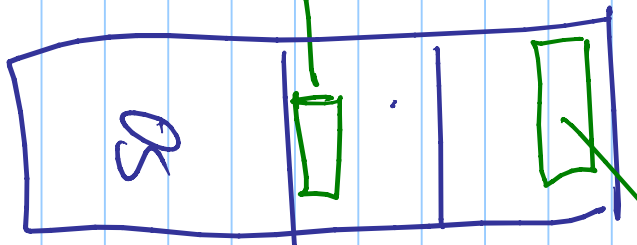
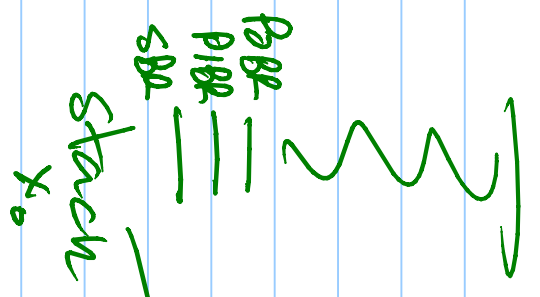
P0 — 0-1GB  
P1 — 1-2GB  
SYS — 2-3GB  
RESVD 3-4GB



Segment table 3 entries



X<sub>0</sub> example - code, data, heap - X<sub>0</sub>

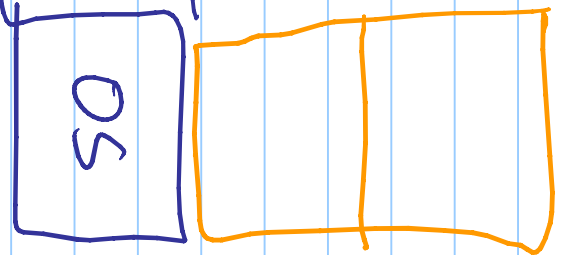


X<sub>0</sub>'s  
 P0 seg  
 P1 seg

sys seg

same

~~X<sub>1</sub>~~



X<sub>1</sub>'s  
 P0 seg

X<sub>1</sub>'s  
 P1 seg

sys

paging systems → pure paging



# Virtual Memory

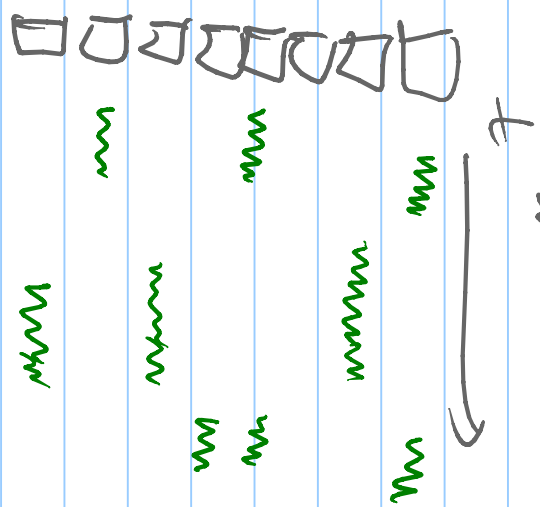
— amount of memory used is larger than available physical RAM

— even 1 process may occupy more mem than all of RAM

Programs have the locality property

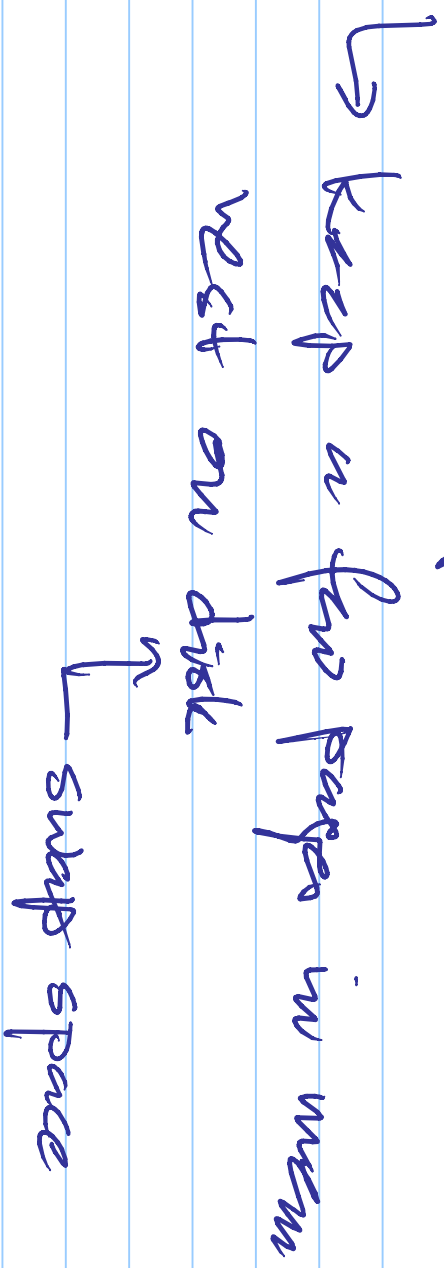
- amount of mem actually used is  $\ll$  total memory allocated

- working set



# Implementing Virtual mem

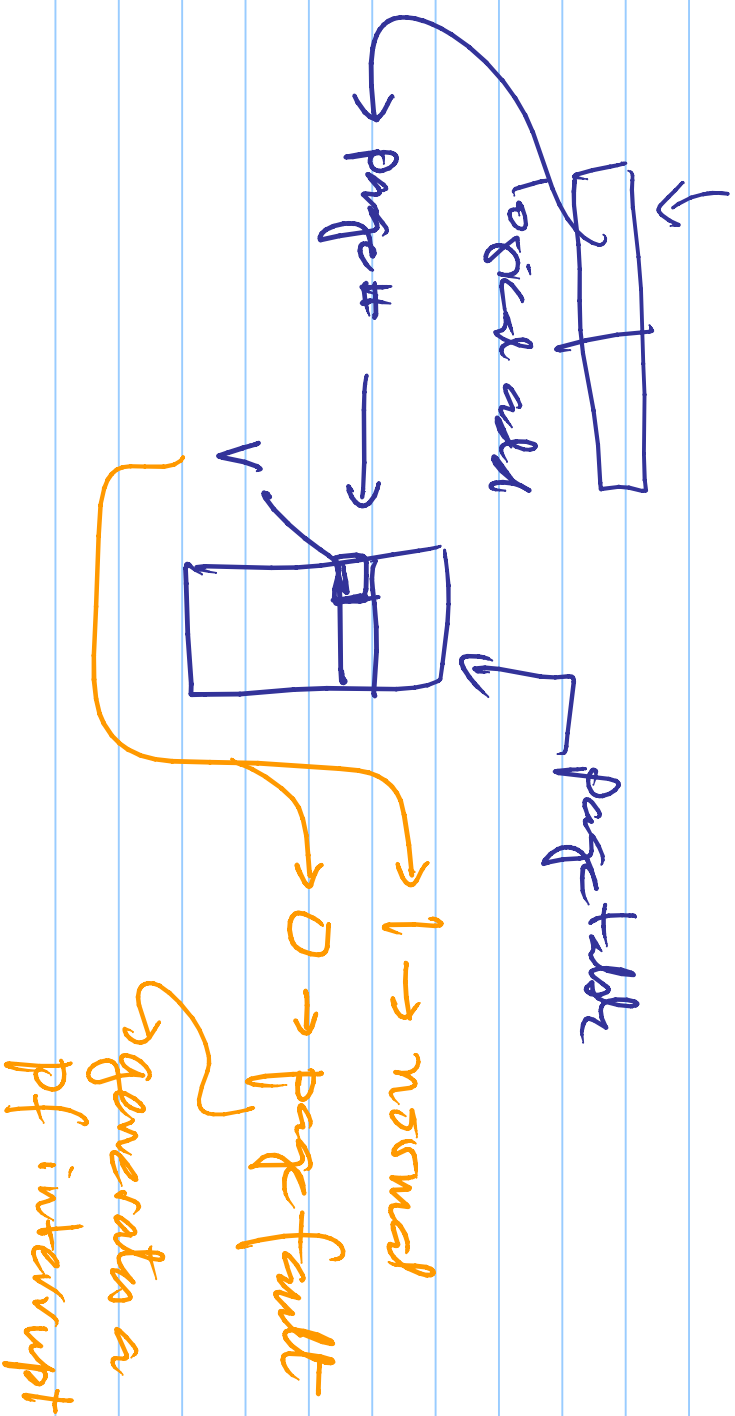
- Several - method
- Demand paging





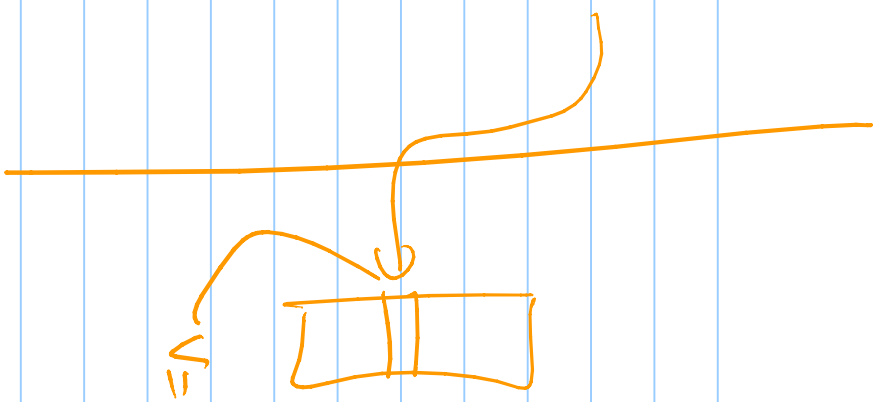
When a process executes

- fetch/load as store



prog

PC → local x  
local y  
⋮  
⋮



PF

→ inst handler

- find disk addr
- find free RAM
- copy disk contents to RAM
- change PTE to have frame addr
- change V bit to 1
- restart instr

Start a process - no pags

working set loaded - some pags

move pf - more pags

→ need to free up mem

↓  
page replacement

Page replacement policies → Page out & Page in

• FIFO

↳ choose victim to page out

• BPT

first page in, first out

• LRU

Best policy

↳ Least recently used

FIFO

→ ① the old page may still be heavily used

- ② suffers from Belady's anomaly
- ③ Not a stack policy

# Page reference string

1 2 3 4 1 2 5 1 2 3 4 5

3 pages  $\rightarrow$  X X X Y Y Z 5 - - 3 4 (9)

4 pages - X Z 3 4 - - 8 1 2 3 4 5 (10)

5 1 - 1 2 3 4 - - 5 - - - (5)

BELADY'S KNOX