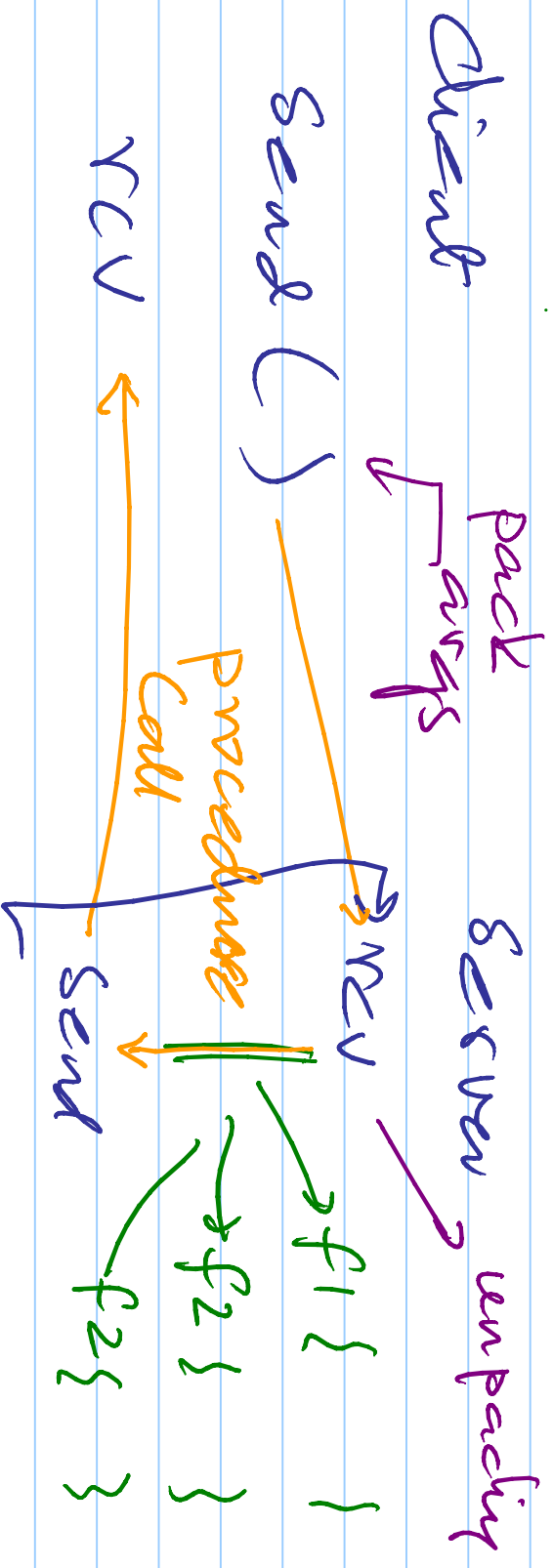


Remote procedure calls (RPC)



RPC systems use a language (IDL)

IDL $\rightarrow$ Interface definition language

Old - DCE-RPC, CORBA, DCOM

New - Java RMI, SOAP, SDA

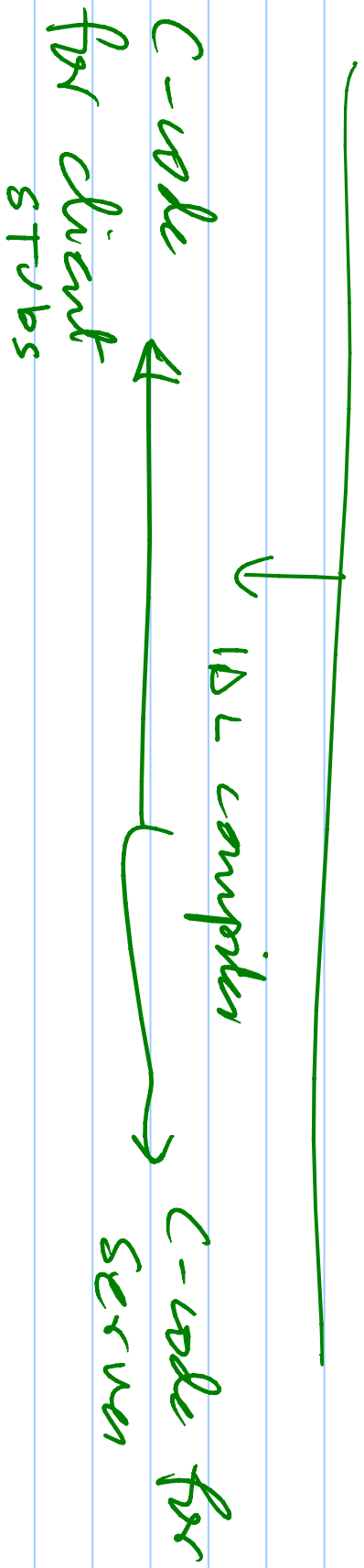
Simple object
access protocol

Service
oriented
architecture

Server code is in IDL

ext function f1 (args) { body }

ext function f2 (arg) { body }



Server $\rightarrow f_1(\text{arg1}) \quad f_2(\text{arg2})$

Client stub

$f_1(\text{arg1})$

+ initialization code
for ports $\checkmark$ OP-code = 0

code to pack args (marshalling)
send (sp, msg) ~~inserts~~ reply_port
recv (reply_port, msg)
unpack args (un-marshall)
}

main()

$f_1(i, j, k)$

Server code

main()

{ registers apart

while loop

rcv(sp, msg)

unpack ms → op-code

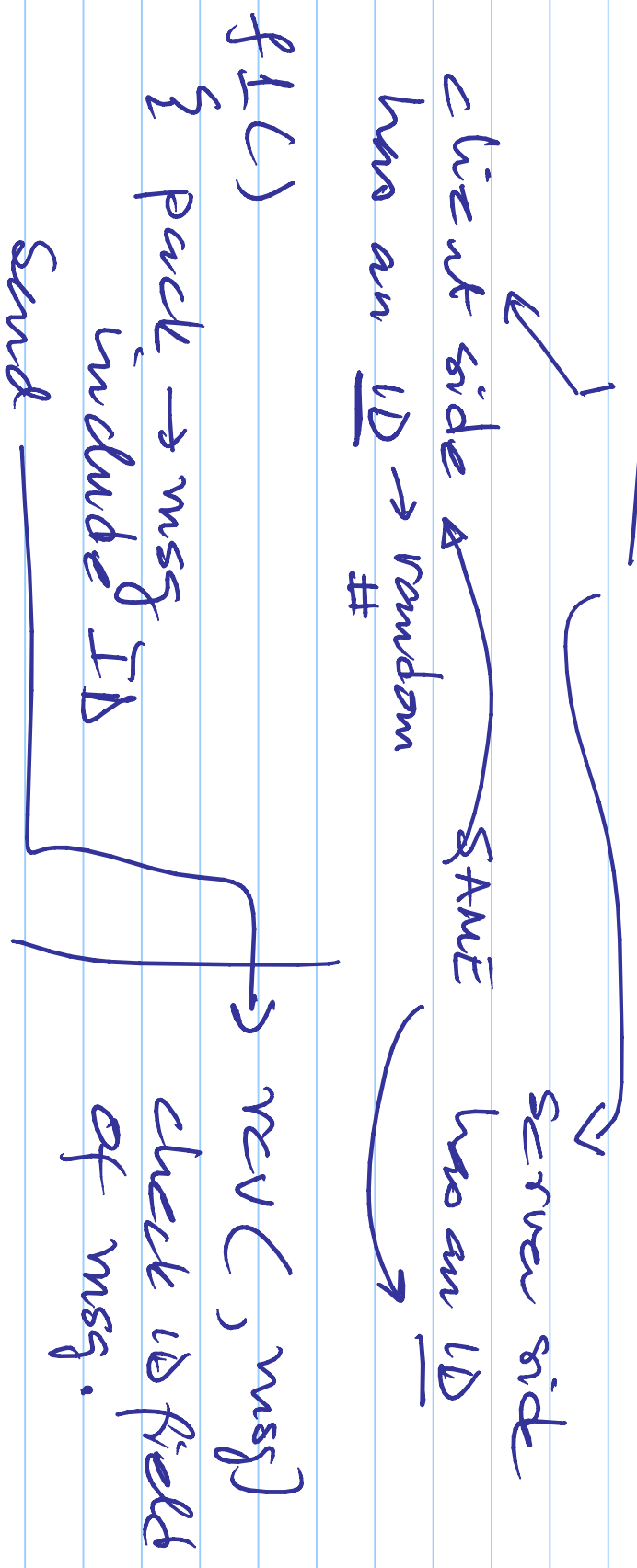
→ further unpack

switch op-code arg

0 → f1(arg1) { code, }

1 → f2(arg2) { code, send }

Version check



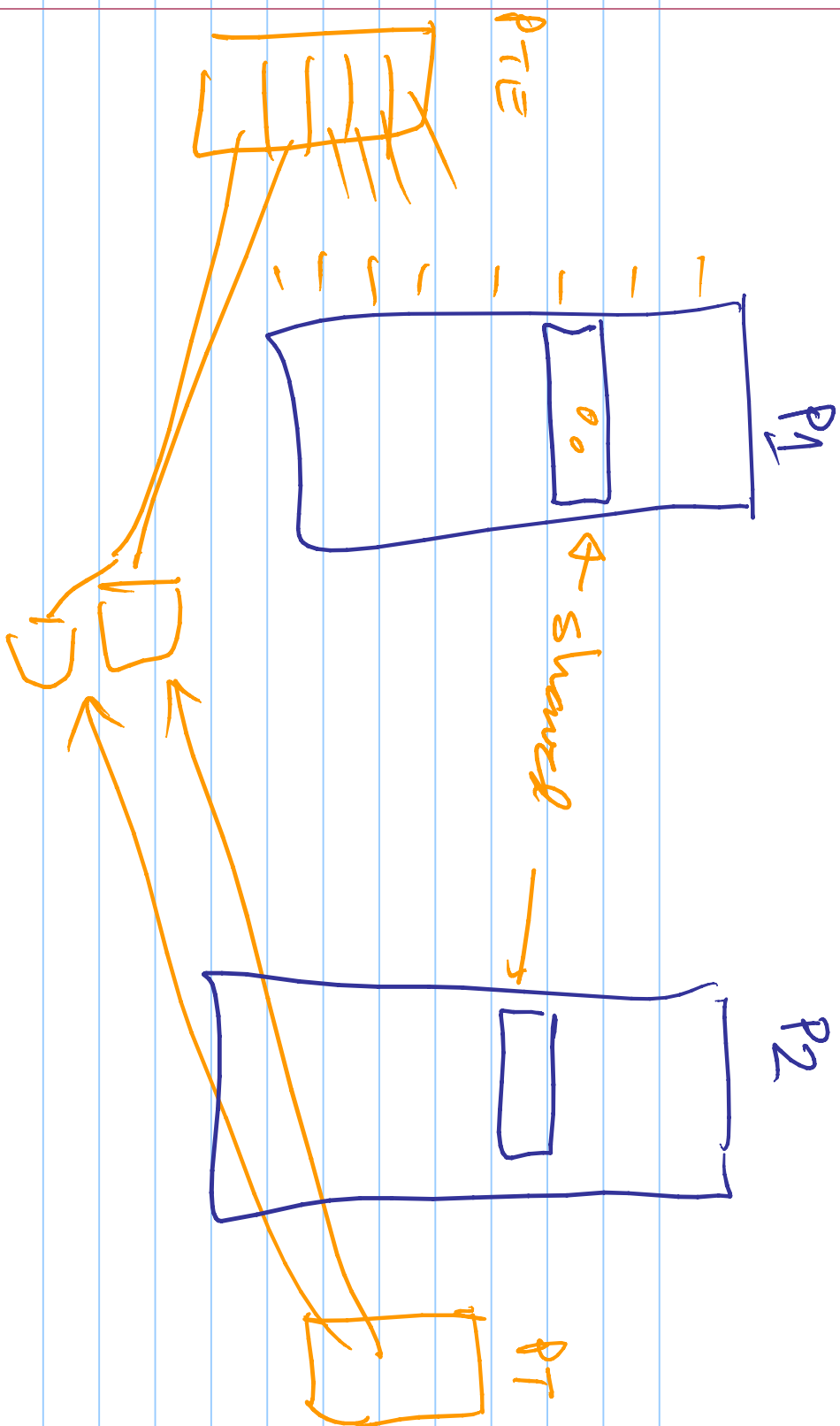
Distributed Shared Memory (DSM)

- Distributed means 'no shared mem'

Shared mem $\rightarrow$ one segment of RAM

$\hookrightarrow$ two or more threads

$\hookrightarrow$ two or more processes



DSM $\rightarrow$ 2 or more processors on
different machines share
one (or more) segments of
memory.

$\swarrow$
 $\searrow$
@RAM level
L0AD/STO mem

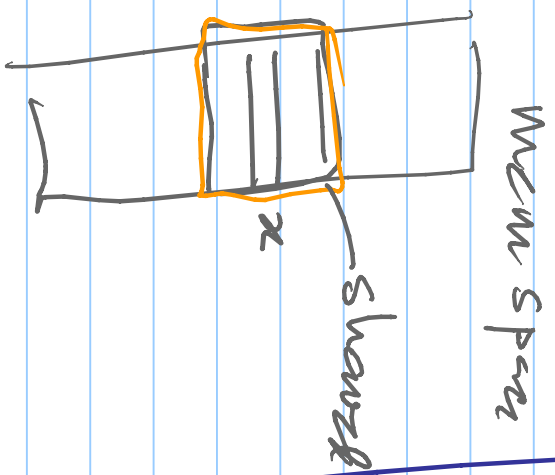
Process $\rightarrow$ DSM client

with x " " "

}

$$x = x + 1$$

}

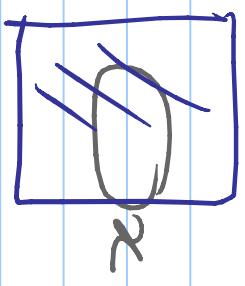


DSM Server

$\rightarrow$ put $\rightarrow$ recu

==

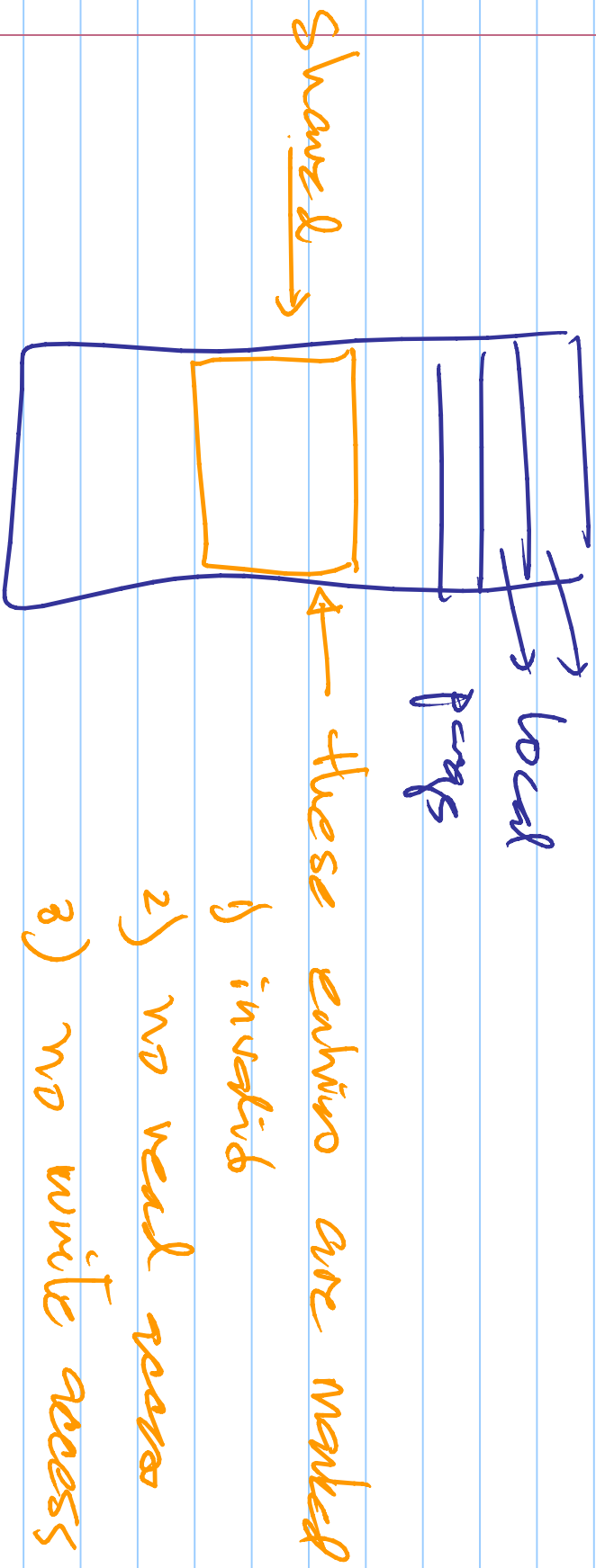
$\leftarrow$ sent



Shared mem

Segment
& contents

DSM client (process) page table



Client read page fault

Read-pf

{ ask server for copy of page
copy read page to mem
update PTE to point to
set valid
set read ok
write not ok
}

↓ for reading

↓
! write get
it
! no locking

Write of handler

for writing

{ is the page in r-mode?

↓

YES

↳ NO → Ask Server for page

get page, install

set w bit to OK

set R bit to OK

}

- Same

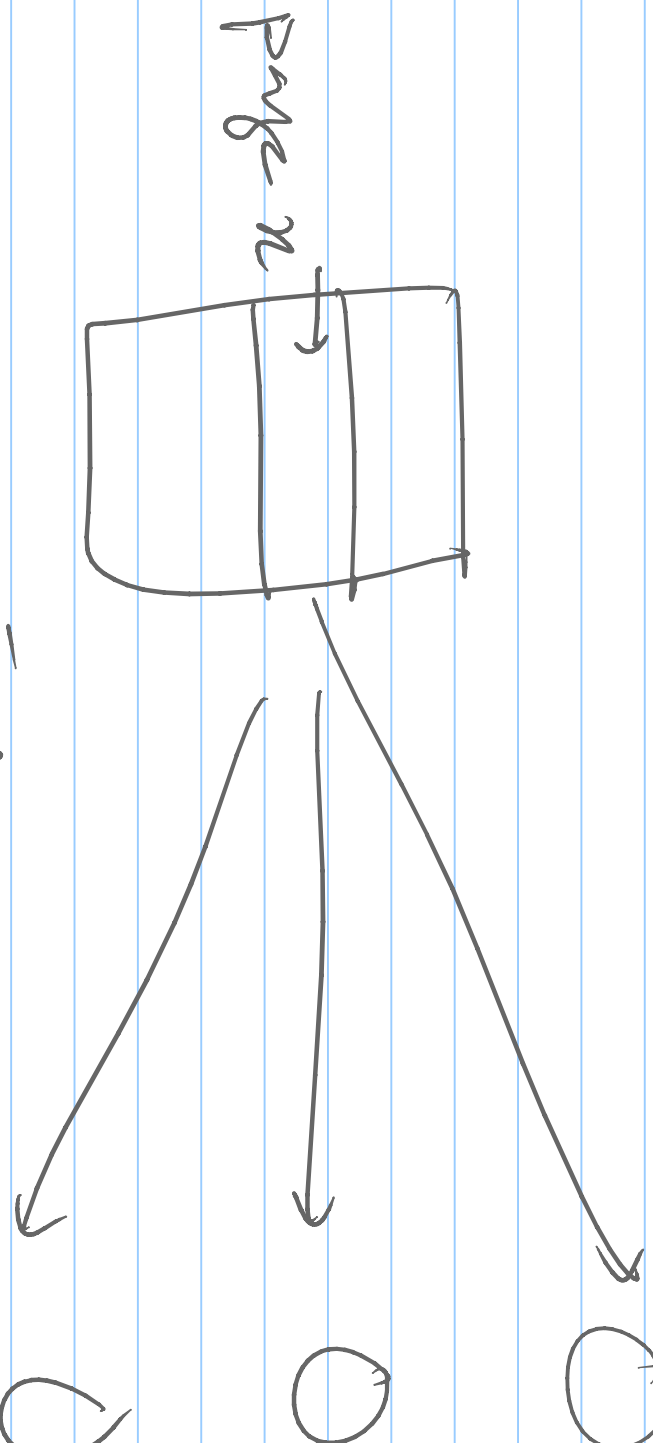
But no need

to get the data

- Server gets request for page in W-mode

Server rules

clients
have it



RULE $\rightarrow$ multiple read copies OR 1 write copy

Invalidation

P_1, P_2 have x in R mode

P_3 gets w-page fault

ask server for w mode on x

server invalidates P_1, P_2

give copy to P_3

Print

1 1 1 1 1 1 1

→ LOAD x

1 1 1 1 1 1 1

STO x

1 1 1 1 1 1 1

Server code

→ x -Page fault (request)

{ if page in R mode or none mode

{ send copy of page,
add client to list of owners
set mode = R }

if Page in W-mode

{ send invalidate to owner

get copy of page (latest)

send copy

add owner

mode = R }

W - Page fault request

{ page in move mode → Send page

Page in R mode

→ Send invalidate to all owners

wait for ack

Send page --- etc

Page in W mode
→ send W-invalidate

get copy of page, Send copy

set mode, owner

invalidation Server at every client

→ r -invalidate

→ set $r, w \rightarrow \text{none}$

→ send ack to server

→ w -invalidate

→ set $r, w \rightarrow \text{none}$

→ send page to server

→ no locking → race condition

→ use locking

→ page shuffling is reduced.

DSM has false sharing

P_1
for $x = x + 1$

P_2
for $y = y + 1$

x & y are on same page!

→ page shuffling happens,

Release Consistent DSM

- no page shuffling
- not coherent

(~~but~~ looks coherent)

- no false sharing